

Implementation and Analysis of Geometric Algorithms for Translucency Sorting in Minecraft

Masterarbeit

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Motivation and Introduction

Translucent Rendering in Sodium



(a) Incorrect rendering

(b) Correct rendering

Figure 1: A Minecraft scene with translucent blocks showcasing the problem.

- \rightarrow This work adds correct translucent rendering to Sodium
- ightarrow Minecraft implements it too, but slower and with errors

- Translucency is common in computer graphics
- Translucency \neq Transparency
- Minecraft/Sodium use ordered rendering of quads with alpha-blending
- Alternative approach: Order-independent transparency
 - Visual inaccuracy
 - Ground-truth variants slower
 - Incompatibility with targeted platforms

- · Implementation of indexed translucent rendering in Sodium
- Improvements to distance sorting
- Special case handling to simplify sorting where possible
- Topological visibility graph sorting
- Plane-based sort triggering
- Partition tree sorting
- Analysis of unaligned partitioning

Rendering and Simple Translucency Sorting

- Blocks are rendered as multiple one-sided quads
- $\cdot\,$ Each 16 \times 16 \times 16 block section is meshed after any change
- \rightarrow Alpha blending: interpolation with alpha factor, non-associative
- → Back-to-front ordering required for correct image
 - \cdot The GPU renders quads in the order of the index buffer

Problem (Quad-based Translucency Sorting)

- Given is a set of **one-sided** quads Q, each consisting of four vertexes $q = (v_1, v_2, v_3, v_4) \in (\mathbb{R}^3)^4$ with $V_q = \{v_1, v_2, v_3, v_4\}$. Each quad's vertexes lie in the same plane $P_q = (n_q, d_q)$ with the unit **normal vector** $n_q \in \mathbb{R}^3$ and **distance** from the origin $d_q \in \mathbb{R}$.
- Quad-based translucency sorting entails finding a **permutation** S of the quads in Q such that quads with greater depth are rendered before those with lower depth.
- A quad q may be **omitted** from S if it is facing away from the camera's position $c \in \mathbb{R}^3$ and thus is not visible.





Figure 3: Intersecting geometry is unsortable.

Figure 2: Cyclically overlapping geometry is unsortable.

 \rightarrow Unsortable geometry requires fragmenting quads or per-pixel sorting



Figure 4: No sorting required if aligned to the bounding box.



Figure 5: No sorting required if opposite facing and only one dot product value.



Figure 6: Quads with **two opposing orientations** can be sorted by dot product. (SNR Sorting)

- Static: One sort order is correct for all view points
- Dynamic: Updated as camera moves and re-sorting is triggered
- Data gathered during meshing determines sort type
- Sort types:
 - None
 - SNR (Static Normal-Relative)
 - Static Graph Sort
 - Dynamic Graph Sort/Partition Tree

Distance Sorting

- $\cdot\,$ Sorting by distance from the camera to the center of each quad
- ightarrow Correct in many cases, but not universally, even with frequent sorting
 - Invalidation is unspecific and is only estimated
 - Used as a fallback when all other methods fail



Figure 7: Center distance sorting is not correct in some situations.

Sorting with Graphs

- $\cdot\,$ Sorting based on potential visual overlap
- → One correct sort order for a set of view points: Camera polytope C
 - Sorting the visibility graph G_C topologically yields a sort order

Definition (Accurate Visibility Condition)

- A quad p is visible through another quad q, referred to as q seeing p, if there is a view ray from a position $s \in C$ to a point $t \in S_p$ that intersects with a point in S_q . S_x is the surface of a quad $x \in Q$.
- Additionally, both quads must be facing the camera to be visible at all with $(s t) \cdot n_p > 0$ and $(s t) \cdot n_q > 0$.

Dynamic Sorting With the Camera Polytope

- ightarrow Camera polytope is constrained by planes of invisible quads
 - Camera exits the camera polytope when a quad becomes visible
 - Triggering is implemented with a combination of interval trees and sorted lists of dot products



Figure 8: A quad q, the camera polytope C in orange, the visible space T_{Cq} in blue (rendered before q), and the intermediary space in green (rendered after q).

Complexity of Visibility Graph Sorting

- The camera polytope has constant complexity with quantized normals
- Evaluating the visibility condition takes constant time
- Visibility is not transitive
- Topological sorting may need to process $O(|Q|^2)$ edges
- ightarrow Finding a visible neighbor efficiently is hard



Figure 9: The arrows indicate the visibility relation, and view rays are represented by dotted lines. The relation is not transitive since q cannot see r.

The Implementation of Graph Sorting and its Limits

- Static sorting for limited size instances with reduced visibility condition
 - Avoids solving a linear program for the accurate visibility condition
 - Assumes visibility from anywhere: $C = \mathbb{R}^3$
 - · Permits testing visibility with just axis-aligned bounding boxes
- For non-static sorting, separators alleviate some cycles
- Finding unaligned separators is infeasible



Figure 10: The separator proves that the blue quad is not visible through the green quad from any point in *C*. A curved view ray is impossible.

Sorting with Partition Trees

Definition (Valid Partition)

- A partition R = (g, d) of \mathbb{R}^3 is a plane described by normal $g \in \mathbb{R}^3$ and distance $d \in \mathbb{R}$.
- A partition is valid, also called a **free cut**, if no quads are intersected by the partition plane.

Definition (Useful Partition)

A partition R is useful if it partitions Q into at least two non-empty subsets.

Partition Trees

- A tree arises from recursive multi-partitioning
- → Indexes are written for partitions in back-to-front order from the camera perspective
 - It always generates a sort order in linear time
 - Sorting is triggered when the camera crosses a partition plane
- \rightarrow Construction in $O(n^2 \log n)$, in $O(n \log^2 n)$ if balanced

Figure 11: The tree degenerates to a list in the worst case.



Figure 12: An example of recursive multi-partitioning on 2D lines.

Multi-Partitioning With Known Orientation

- 1. Projects the quads along a known orientation
- 2. Sorts the interval endpoints
- 3. Scans the list for gaps that permit partitions



Figure 13: Projection along two axes yields two gaps for each one.

- ightarrow The orientation is known: Axis-aligned partitioning only
 - Generates as many partition planes as possible
 - \cdot Quads can be on the partition plane

Unpartitionable Instances





Figure 15: These unaligned quads are not partitionable.

Figure 14: These aligned quads are not partitionable, even when certain quads are removed.

- Detecting special cases during partitioning allows for fewer nodes and partition planes
- · Partial tree updates improve partitioning speed for large sections
- · Index compression reduces memory usage very slightly
- · Primary intersector detection avoids failure on intersecting geometry

- Partition tree sorting is faster than distance sorting by 60%
- · Static visibility graph sorting avoids some dynamic sorting
- Detecting static special cases is important
- Distribution of sort types varies significantly with scene content
- \rightarrow Distance sorting is never used in practice

Unaligned Partitioning

	non-intersecting		
	sortable		
unaligned	unaligned partitionable		
intersecting	aligned partitionable		aligned
unsontable			
	SNR sortable		
	always sorted	᠕	
		IJ	

Problem (Unaligned Partitioning (UAP))

Given a set of n partitionable quads Q, does a valid and useful partition R exist, and what are its parameters?

- Only known-orientation partitioning is fast
- \rightarrow Large solution space of unknown-orientation partitions
 - No hard cases of UAP naturally appear in Minecraft



Figure 16: Aligned geometry can require an unaligned partition.

- Slope-offset parametrization $y \ge ux + vz + d$ of planes yields parameter space $\mathbb{L} = \mathbb{R}^3$
- → Linear mapping from $p \in \mathbb{R}^3$ to the parameter-space half-space containing all $R \in \mathbb{L}$ for which p is in the real-space half-space $H_{\mathbb{R}}^2$
- → Parameter-space faces correspond to real-space vertexes
 - Valid partition sets and useful partition sets as expressions over parameter-space half-spaces



Figure 17: Valid partition set highlighted in red for blue quad *q*.

Useful Partitions in Parameter Space



Figure 18: Useful partition set highlighted in red for the three quads.

- Combinatorial solution in $\Theta(n^4)$
 - 1. Pick all triples of vertexes to form a plane
 - 2. Test the partition for validity and usefulness
 - 3. Mitigate a global useless autopartition with an artificial vertex
- Incremental solution with **linear constraint sets (LCS)** possible, but worse time complexity
- Can also be solved with LCS-SAT

- $\cdot\,$ Depending on the input constraints to UAP, it can solve LCS
- Transfer of lower bound on LCS to UAP
- ightarrow Maps half-space constraints to vertices
 - $\cdot\,$ Requires more quads and expensive collision avoidance if UAP is strict

Conclusion

- Usually translucency requires OIT
- Minecraft's specific type of geometry makes quad-based translucency sorting tractable
- ightarrow Taking advantage of many special cases and well partitionable geometry
 - \cdot Ground-truth results are possible, faster than distance sorting
 - Correct translucent rendering without hardware support
 - Unaligned partitioning is infeasible, but polynomial

- Is an acceleration structure for the reduced visibility condition feasible and useful?
- Can visibility graph sorting within the partition tree improve its characteristics?
- To what extent can convex partitioning, a superset of unaligned partitioning, form a bridge between partition and graph-based sorting?

Thank you for listening!